# THE WONDERFUL FAIRY OPTIONS OF WinChloe

by IGM Petko Petkov

#### Dedicated to Christian Poisson

#### I. FOREWORD

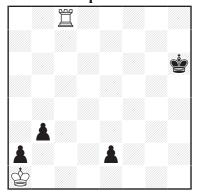
I write this article because of a very interesting new phenomenon in the world of fairy composition: the special options **Conditionsfromply=n** and **Conditionsuntilply=n**, invented by the world-known programmer Christian Poisson and practically demonstrated in version 3.36 of his wonderful program WinChloe.

More than a year has passed since the discovery of Christian Poisson, but to my surprise there are still few published problems of this type. But I think that WinChloe's options open up very rich and interesting possibilities for the composition of fairy problems of the hypermodern form! Therefore, I decided to write this article with the hope that it will be of interest for composers and solvers all over the world! In fact, the beginning of this very thrilling story happened with the traditional 12<sup>th</sup> Belgrade Problem Chess Festival 2017 (26<sup>th</sup>-28<sup>th</sup> May 2017). The Serbian organisers, who have always been authors of very interesting events and competitions, this time offered in the Belgrade Internet Tourney 2017 the following very fresh, non-standard theme:

Group C – hs# maximummer 2,5 & 2 solutions; Thematic condition: Help-selfmate Maximummer with help-play before the last two half-moves, presenting s#1 Maximummer (Black has to play the geometrically longest move). No other fairy conditions or pieces are allowed.

The theme was illustrated by the following example (note that the number of solutions must be exactly two, and the number of moves is also strictly defined - 2,5 moves!):

#### 1 - Example - Scheme



hs#3 0.2.1.1.1.1 (2+4) C+ Maximummer

#### Solutions:

I. 1...é1= 2. ☐ f8 2 h4 and now Black are required to play using the condition Maximummer: 3. ☐ f6+ the finale is the geometrically longest move by Black: 3... 2×f6#

II. 1...é1= \(\boxed{\boxes}\) + 2. \(\boxed{\boxes}\) \(\boxed{\boxes}\) h1 and then follows: 3. \(\boxed{\boxes}\) b1 \(\boxed{\boxes}\) × b1#

We will demonstrate other examples from the thematic tournament, which turned out to be quite interesting.

After the tournament was announced, I asked WinChloe's author, Christian Poisson, whether it was possible to create a computer program to solve problems of type hs#2,5 Maximummer? I was extremely surprised when Christian wrote to me a day later that he had already created a program that had much more potential and after a few days of testing he would include it in the new version of WinChloe, at that time 3.36!

Above all, I want to express my great gratitude and respect to the world-known author of WinChloe Christian Poisson! My suggestion is that this kind of fairy problem, with WinChloe's special options, be named TYPE POISSON!

Of course, Christian's new program can also solve the hs#2,5 Maximummer problems in just a few seconds! After the new version of 3.36 (in May 2017) was released, I immediately announced this in the forum on the world-known Serbian site **MatPlus**, for the benefit of all participants in the 12<sup>th</sup> Belgrade T.T. (section C - Fairies).

Marjan Kovačević and Borislav Gadjanski have made a great contribution by organising this tournament - two great names in the world of chess composition that do not need a detailed presentation! The first journal to publish Poisson-type problems was the United States StrateGems - in its issue n°79 (July-September 2017). Later, new originals of this type appeared also in this magazine. The Bulgarian site **KoBulChess** also published in 2017 the detailed definition and explanations of these options and an original hs# problem by P.A. Petkov.

#### II. FEATURES OF THE NEW OPTIONS IN WinChloe

In brief, the essence of the program was explained by its inventor - Christian Poisson, in «Readme» of the new (in May 2017), version 3.36 :

«Two new options are programmed, allowing to indicate from which ply, or until which ply, fairy conditions are taken into account. These options, to write in the «twin» field, are:

**Conditionsfromply=n** (or in French: Conditionsapartirde=n) and

**Conditionsuntilply=n** (or in French: Conditionsjusqua=n), «n» being the number of plies.

For example: Conditionsfromply=5 means that the conditions are taken into account only from the  $5^{th}$  ply (the  $3^{rd}$  white move for a direct, a self or a help-self problem) and the  $3^{rd}$  black move for a help-problem. Conditionsuntilply=6 means that the conditions are taken into account only until the  $6^{th}$  ply (the  $3^{rd}$  black move for a direct, a self or a help-self problem and the  $3^{rd}$  white move for a help-problem). It is possible to combine both options (write them on to different lines).

IMPORTANT: THESE OPTIONS WORK WITH «SIMPLE» CONDITIONS (Maximummer, Minimummer, Circe, Sentinelles...), BUT CAN (MAYBE) GIVE UNEXPECTED RESULTS WITH CONDITIONS AS-KING MORE COMPLICATED CALCULATIONS. THEY DON'T WORK WITH CONDITIONS WHOSE EFFECTS ARE NOT CALCULATED AT EACH MOVE.

For more information you can read about the version 3,36 (the Readme file, named "readme.txt" is in the "C:\WinChloe" folder).

Of course, some additional explanations are needed here, which I wrote with the help and control of Christian, who on this occasion answered me very kindly and comprehensively on my long letter sent to him on May 6<sup>th</sup>, 2017.

What does the word «ply» mean? The answer is elementary - «ply» is a half-move. Consequently, each problem has a certain number of «X» «plies» (half-moves) = sum from all single half-moves by White and Black. Examples:

#### a) Orthodox problems and direct non-orthodox problems:

A #2 has X=3 plies, they are:  $1^{st}$  half-move (key move) by White  $+ 1^{st}$  half-move by Black (defence)  $+ 2^{nd}$  move by White (mating move). A #3 has X=5 plies, a #4 has X=7 plies, a #5 has X=9 plies and so on. Here it is not difficult to define a simple formula: if in an orthodox problem the condition is #n, the sum «X» of plies (half-moves) = 2n-1. For example, a #11 has  $2\times11-1=21$  plies. In the general case, this formula also applies to unorthodox direct-problems of the form #n, =n, s#n, s=n r#n, r=n etc.

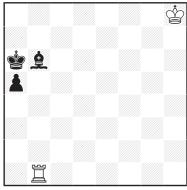
### b) Helpmates (help stalemates, help double-stalemates, help-checks, help castlings, etc) and also Help-selfmates (hs#, hs=, hs==, etc):

A helpmate in «n» moves where «n» is a whole number (an integer) has  $X=n\times 2$  half-moves (plies). For example, a h#2 has  $X=2\times 2=4$  plies, they are :  $1^{st}$  move by Black +  $1^{st}$  move by White +  $2^{nd}$  move by Black +  $2^{nd}$  (mating) move by White ; a h#3 has  $3\times 2=6$  plies, a h#4 has 8 plies etc...

Very Important! If «n» is a fractional number, the number of plies is also X=2n+1 half-moves! For example, h#1,5 has  $X=2\times1,5+1=4$  half-moves (plies) H#2,5 has  $X=2,5\times2+1=6$  (plies), H#3,5 has 8 plies etc. The same situation occurs with hs#n problems in which «n» is also a fractional number, etc!

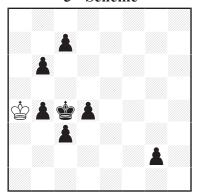
At first sight these words appear to contain an error. The logical question can be: if we have for example a h#1,5 this problem has de facto only 3 plies: the first move by White + the second half-move by Black + the last-third half-mating move by White. BUT! By using WinChloe's options here, we must conditionally assume that such problems also have 4 plies but the first (missing) half-move of Black is assumed as a «null ply». In other words, we must know that according to WinChloe the play begins here with the second ply (de facto - the first half-move of White) followed by the 3<sup>rd</sup> ply (the second half-move by Black) and the end follows the 4<sup>th</sup> ply (the third mating half-move by White). These rules have a very important practical value and they should be applied accurately! Errors in this respect lead to erroneous results by the program. Let's see a simple example, scheme n°2:

#### 2 - Scheme



h#2 0.1.1.1. (2+3) C+ ConditionsFromPly = 4 Sentinels

#### 3 - Scheme



hs#3 0.1.1.1.1.1 (1+7) C+ ConditionsUntilPly = 2 Anti-Andernach

n°2. Here it is clear that the condition Sentinels must be applied only on the last half-move - the mating move of White. But since we conditionally accept that this problem also has 4 plies, with the first ply being the absent first half-move of Black («null ply»), the play starts here on the 2<sup>nd</sup> and ends on the 4<sup>th</sup> ply! For this reason, we must write under the diagram ConditionsFromPly=4.

Solution: 1... \( \begin{aligned} \begin{alig

But if we write erroneously ConditionsFromPly=3, WinChloe claims that the problem does not have a solution! The reason is clear - the program perceives here as a third half-move (ply) the black move 2.247(+1.06) and the Sentinels condition is mistakenly applied (before we need it). As a result, a black 1.06 bb closes the line for the White rook and 1.06 bb is an impossible move!

n°3. This is an analogous example with a hs#2,5. Here the condition Anti-Andernach is applied only on the first (key-move) by Black but WinChloe perceives this move as a second ply (the first missing half-move by White is a «null ply») and under the diagram we must write ConditionsUntilPly=2.

Solution: 1...g1=\( \mathbb{U}(\mathbb{U}) 2.\mathbb{U}g5 \( \cdot \) 6 3.\( \mathbb{U} \) b5+ \( \cdot \times \) b5#

Of course, here too it is erroneous to write ConditionsUntilPly=1 - WinChloe does not give a solution!

#### c) Series-movers (help-play) - Ser.H#n, Ser.S#n, Ser.R#n, etc.

Here, in a series-problem in «n» moves the number X of all plies is n+1. In other words, we have n-1 half-moves (plies) and only the last n<sup>th</sup> move is a compound move: it includes the last (n<sup>th</sup>) half-move (ply) of the party carrying out the series and the n+1<sup>th</sup> last half-move of the other party, which obliges the mate, the stalemate, etc...

#### d) Series- Autostalemate in «n» moves (ser-!=n)

The number X of all plies is also n.

The basic principle of applying WinChloe's options is very interesting: the play in each problem with X plies in the general case contains two parts:

#### A) Orthodox part, where the number of plies = A

Here the play is carried out only according to orthodox rules. Very important: in each problem the presence of the orthodox part A is obligatory! In other words: a problem with conditions from ply (or conditionsuntilply) cannot exist without the orthodox part A!

#### B) Fairy part - here the play containing B plies

He is carried out according to the rules of one or more (two, three, four, etc...) fairy conditions, for example: B= Circe, or B= Circe + Maximummer, etc...

Consequently, in accordance with the structure of these problems, the common sum X of the half-moves (plies) can be:

$$X = A + B \text{ or } X = B + A$$

In other words, the play can begin with the orthodox part A, followed by the fairy part B (Option = ConditionsFrom Ply) or vice versa (Option = ConditionsUntil Ply).

A more complex variant is possible using both options as we shall see later.

#### 1. Use of option CounditionsFromPly

As we already noted above, the first part here has A half-moves (plies). Then follows the second fairy part of B half-moves (plies).

For example : a h#3 has 6 plies : 1st ply = first half-move by Black, 2nd ply = first half-move by White, 3<sup>rd</sup> ply = second half-move by Black, 4<sup>th</sup> ply = second half-move by White, 5<sup>th</sup> ply = third half-move by Black and  $6^{th}$  ply = third white mating move. The composer himself chooses a fairy condition (or combination of fairy-conditions) for the fairy part of his problem. After that, he chooses from what half-move (ply) the use of this fairy-condition commences. If we have 4 orthodox plies (A = 4) plus two fairy-plies (B = 2) then this means that at first we make 4 orthodox half-moves and only with the 5th ply does the execution of the fairy condition commence.

In this case we must write : ConditionsFromPly = 5.

Solution of  $n^{\circ}4$ :

1. **■** g5! **\$** ç4 2. **■** gb5! **\$** a2

and then follows : 3. **a** a4(+ **a** a3) **a** b3(+ **a** a2)#

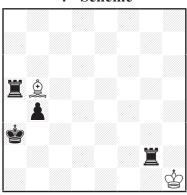
Returning to the scheme n°1, we can note that here we can write the stipulation under the diagram in a new way:

hs#2,5 & ConditionsFromPly=5 Maximummer. After solving this problem by computer, you need to write below the diagram: 2 solutions:

I. 1...é1= **≜** 2. **□** f8 **≜** h4 and here, from the 5<sup>th</sup> ply, we must apply the condition Maximummer: 3. \( \begin{aligned} \frac{1}{6} + & \text{ followed by the last (mating) move} \end{aligned} \)  $(6^{th} ply): 3... \stackrel{\bullet}{\cancel{2}} \times \mathbf{f6}\#$ ; analogically:

II. 1...é1= $\mathbf{Z}$ + 2. $\mathbf{Z}$ ¢1  $\mathbf{Z}$ h1 and here follows : 3. $\mathbf{Z}$ b1  $\mathbf{Z}$ ×b1# Under the diagram of each problem from Belgrade Internet Tourney (BIT) 2017 Fairy T.T., one can write exactly the same text as seen in scheme n°5.



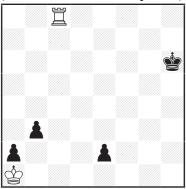


h#3 (2+4) C+ConditionsFromPly = 5

Sentinels

#### 5 - Scheme

(as 1 but with list of options)



hs#3 0.2.1.1.1.1 (2+4) C+ ConditionsFromPly = 5Maximummer

We stress that the option ConditionsFromPly = 5 should be written in the «Twin» field of the program WinChloe, but the condition Maximummer must be marked in the field «Conditions» using the ordinary practice of such a choice!

It is clear that under the diagram of Scheme  $n^{\circ}6$  we can also write down the option ConditionsFromPly = 6, since de facto here the 5<sup>th</sup> ply is also an orthodox half-move by White (the condition Maximummer applies only to Black!).

Demonstrating the enormous possibilities of the WinChloe's options, first of all I want to show several of the most interesting problems from the T.T. Belgrade Internet Ty 2017.

#### Solutions of n°6:

I. 1... **≝**g1 2. **豐**×g1 **≜**f8 3. **豐**g8 **豐**é7#

II. 1... 宣f1 2. 豐×f1 宣g4 3. 豐h3 豐f4#

Comment by Emanuel Navon in *Variantim* (n°72/2017): «A complex combination of sacrifice and pinning (of attacking pieces), two characteristic elements that prevent longer moves by Black. The whole content is harmonized in ODT form, with Bristol clearance, black halfpin, and anticipatory self-pins.»

Solutions of n°7:

I. 1... ♀ h4 2. ∅ é3 d8! and then follows the s#1 Maximummer:

3. \( \begin{aligned}
2 & \delta & \de

Analogically:

II. 1... 置h3 2. 分f6 豐b3! 3. 罩b1 置ç3#

Solutions of n°8:

II. 1... 豐×b5! 2. 罩g4 豐a6!! 3. 罩é4! 豐×f6#

Extreme activity of the black Queen which performs all six(!) black half-moves! Here are particularly interesting the annihilation captures of  $\Xi g7$  and  $\Xi b5$  that lead to pins of the white Knights, the blocks realized by the white Rooks on é4 and f4, the reciprocal change of functions in both solutions. The finales are beautiful model pin-mates, realised in Meredith form.

Solutions of n°9:

I.1... 罩×ç2 2. ②×f4 皇 ç6! 3. 罩 d3 豐×f4#

II. 1... **≜** ×f3 2. **⊴** ×b2 **≡** c6! 3. **≜** d3 **⋓**×b2#

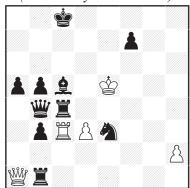
A very non-standard and beautiful demonstration of the Black Grimshaw theme after critical moves of **Z**-**Q**!

This beautiful idea is complemented by the sacrifices by white Knight, white blocks on d3 and the surprising mates by the black Queen.

#### 6 - P. Einat

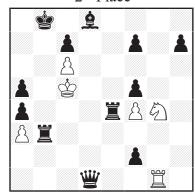
Belgrade Internet Ty. 2017 (v) 1st Place

(Version by P.A. Petkov)



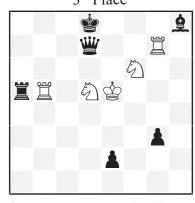
hs#3 0.2.1.1.1.1. (5+10) C+ ConditionsFromPly = 5 Maximummer

## **7 - M. Klasinc**Belgrade Internet Ty. 2017 2<sup>nd</sup> Place



hs#3 0.2.1.1.1.1. (6+12) C+ ConditionsFromPly = 5 Maximummer

### **8 - P. Petkov**Belgrade Internet Ty. 2017 3<sup>rd</sup> Place

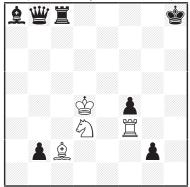


hs#3 0.2.1.1.1.1 (5+6) C+ ConditionsFromPly = 5 Maximummer

#### 9 - A. Bulavka

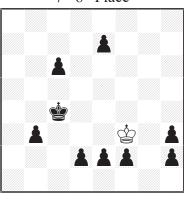
Belgrade Internet Ty. 2017 (v) 4<sup>th</sup> Place

(Version by M. Klasinc)



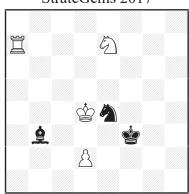
hs#3 0.2.1.1.1.1 (4+7) C+ ConditionsFromPly = 5 Maximummer

#### **10 - I. Tominić** Belgrade Internet Ty. 2017 7<sup>th</sup>-8<sup>th</sup> Place



hs#3 0.2.1.1.1.1 (1+9) C+ ConditionsFromPly = 5 Maximummer

#### 11 - G. Foster StrateGems 2017



hs#4 2.1.1... (4+3) C+ ConditionsFromPly = 7 Maximummer

#### Solutions of n°10:

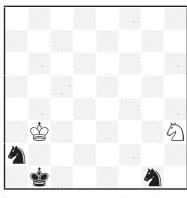
I. 1...h1=খ+! 2. 當 é3 é1= 置+! 3. 當×d2 খd5#

II. 1...é1=**≜**! 2.**Ġ**é2 h1=**△**! 3.**Ġ**f1 **△**g3#

A real sensation - White Rex Solus + black AUW! This is probably one of the few possible positions of this kind! On the other hand, it is obvious that there are no further strategic elements. But nevertheless the author deserves praise! It is necessary to note that so far only a small number of problems of WinChloe's options have been published. The world-known American magazine StrateGems was the first that very quickly (two months after the publication of the results from Belgrade Internet Ty 2017) began popularising this novelty from issue n°79/2017.

Here follows a small series of problems of WinChloe's options published in StrateGems. We start with two beautiful miniatures by the Australian maestro Geoff Foster.

#### 12 - G. Foster StrateGems 2017



hs#5 2.1.1... (2+3) C+ ConditionsFromPly = 9 Alphabetic

#### Solutions of n°11:

I.1. 公 c6 皇d1 2. 曾d5 曾f4 3.d4 皇h5 and follow: 4. 單f7+ 皇×f7#! II.1. 單a4 曾g3 2. 曾é3 如d6 3. 罩é4 皇c4 and follow: 4. 公f5+ 如×f5#

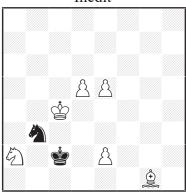
In a miniature form is realized a very nice complex: ideal mates, reciprocal change of functions between two duets of pieces: 2a7/667 and 3b3/664, self-blocks.

#### Solutions of n°12:

I. 1. ∅ f2 ♠b4 2. ὑ ç3 ♠ é2+ 3. ὑ d2 ♠g3 4. ὑ é1 ὑ ç1 5. ∅ d3+ ♠ ×d3# II. 1. ∅ f4 ὑ a1 2. ὑ ç2 ♠f3 3. ὑ d1 ὑ b2 4. ∅ é2 ὑ b1 5. ∅ ç3+ ♠ ×ç3#

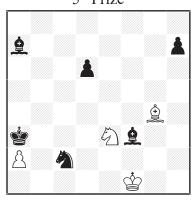
Another interesting Tanagra problem (only 5 pieces on the board!) in aristocratic form and with chameleon-echo ideal-mates! The use of the fairy condition Alphabetic Chess seems very fruitful but so far this opportunity has been demonstrated only in two problems. The second opus is the following hs#7.

13 - G. Foster Inédit



hs#7 (6+2) ConditionsFromPly = 6 Alphabetic

14 - P. Petkov StrateGems 2017 5<sup>th</sup> Prize



hs#4 2.1.1... (4+6) C+ ConditionsFromPly = 7 Minimummer

Solutions of n°13: 1.é4 **a**a1 2.**a**d4 **a**d2! 3.**a**é3+ **a**é2! 4.**a**c1+ **a**d1 6.**a**b2+ **a**c2 7.**a**c2 **a**b3# Round-trip **a**c2-d2-é2-é1-d1-ç2 and switchback **a**b3-a1-b3. A surprising and beautiful finish with model mate!

In the following problem, the condition Minimummer is applied for the first time. The play here also seems very non-standard. Solutions of  $n^{\circ}14$ :

I 1. \( \hat{2}\) f5! d5 2. \( \hat{2}\) ×h7! d4 3. \( \hat{2}\) g8 d×é3 4. \( \hat{2}\) b3! é2#
II. 1. \( \hat{2}\) f5! h5 2. \( \hat{2}\) ×d6! h×g4 3. \( \hat{2}\) é4 g3 4. \( \hat{2}\) c3! g2#

Here the reciprocal change of functions is realised in the play of two duets of figures: 2g4-4d6 and 4h7-4d6. White captures one of the enemy Pawns to prepare the final mate by the other. This annihilation + Zilahi is combined with taking away a field from the black King. Model-mates in Meredith form.

Here follow two very interesting problems by the world-known American composer Kostas Prentos - the editor of the «Retro and PG» section in the American magazine *StrateGems*.

Solutions of n°15:

- a) 1... 🕯 ×ç4! 2. 🖄 d3! 🕭 f7 3. 🕸 b5 👲 g6 4. 🕸 ç4 🕭 ×d3#
- b) 1... 🖺 ×d4! 2. 🖄 d7! 🖺 h4 3. 🕸 d5 🖺 h7 4. 🕸 d4 🗒 ×d7#

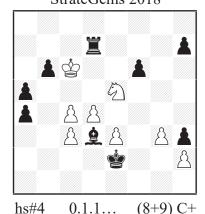
Double realisation of the «Rundlauf» theme, reciprocal change of functions between **Z-2** and of course, a very interesting demonstration of the Umnov theme after the key moves:

In a) 1... **≜**×ç4! 2. ∅d3! and in b) 1... **Ξ**×d4! 2. ∅d7! Excellent!

Solutions of n°16:

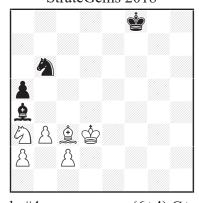
a) 1. ½ b2 ½×b3 2.ç3 ½ f7 (2... ½ é6?) 3. ½ ç2 ½ é8 4. ½ b3 ½ a4# b) 1. ½ é1 ½×a3 2. ½ ç3 ½ é7 (2... ½ d6?) 3. ½ b2 ½ g5 4. ½ a3 ½ ç1# A similar theme to n°12 but here the double «Rundlauf» is performed only by the black Bishop, the finales are model mates and the form is Meredith.

15 - K. Prentos StrateGems 2018



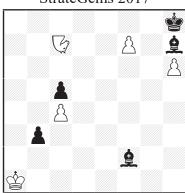
hs#4 0.1.1... (8+9) C b)  $\blacktriangle$  h7 $\rightarrow$ g6 ConditionsFromPly = 7 Maximummer

16 - K. Prentos StrateGems 2018



hs#4 (6+4) C+ b) **2** a4→ç1 ConditionsFromPly = 7 Maximummer

**17 - P. Petkov** StrateGems 2017



hs#3 0.1.1.1.1.1 (5+5) C+ ConditionsFromPly = 5 a) Maximummer b) Minimummer

This light problem, n°17, demonstrates a nonstandard opportunity in this arena - a twin with the replacement of the fairy-condition :

- a) 1... \( \Delta b1! 2. \( \Q \) f1 \( \Delta a2 3. \( \Q \) g3 \( \Delta \) d4#
- b) 1... \( \) \( \

The following Tanagra problem is the first one in which the fairy part B contains two fairy conditions.

n°18. We assume conditionally that the absent 1<sup>st</sup> half-move of White is a «zero-ply». Therefore here the orthodox part A has 2 half-moves (plies) - the 1<sup>st</sup> black half-move as a 2<sup>nd</sup> ply and the 2<sup>nd</sup> white move as a 3<sup>rd</sup> ply. But from the 4<sup>th</sup> ply the fairy part B starts, which is a combination of two fairy conditions: Disparate and Anti-Andernach!

I. 1... ■ b6 2.ç4 end of the Orthodox part A! Now follows the fairy-part B: 2... ₩d7(₩)!! This is not a self check by Black because the white Queen is paralysed for White! 3.ç5(▲)+!! This is a check to the black King from the unparalysed white Queen plus black block on ç5! 3... ★ ×d7#! This is mate to the white King because the black King paralyses the white King! Similarly:

#### II. 1... **ὑ** ç5 2. ❖ h5 d6(쌀)!! 3.ç4(▲)+!! ❖×d6#!

Chameleon-echo model mates after rich strategic play.

I have no doubt that every reader, looking at problems n°1-14, will be surprised at the fact that in almost all cases (with little exceptions) the composers have used only the stipulation hs#. The explanation is simple: this is probably so because the first fairy T.T. of Belgrade Internet Ty 2017 had the same stipulation and then probably «by some inertia» many authors began, also only with hs#, using conditions Minimummer or Maximummer.

Of course, we can use WinChloe's options with all other stipulations but so far there have not been any such published problems.

#### 2. Use of the option ConditionsUntilPly=n

Here we have : a fairy half-moves (plies) + B orthodox half-moves (plies).

As far as I know, the first problem of this type by me was published on September 17<sup>th</sup> 2017 in the famous Bulgarian problem site *KoBul-Chess*. The problem is dedicated to Christian Poisson! We will analyse this problem in detail, but first of all let us consider some easy examples of this type.

n°19. Here option ConditionsUntilPly = 1 means that condition Anti-Andernach is applied only on the first ply (white key-move). After this, two orthodox half-moves of Black and White follow.

Solution: 1.b8=∅(♠)! blocus

The end of fairy condition is followed by two defences by a black Springer:

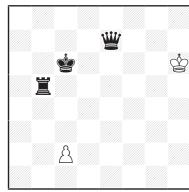
1...**♠**ç6 2.**②**ab7#

1... **△** d7 2. **△** cb7#

A small thematic try : 1.b8 = 2 (2)? blocus

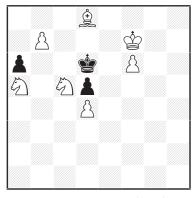
1...**≜** ç7 2.**≜** é7# but 1...**≜** a7!

**18 - P. Petkov** StrateGems 2017 Commendation



hs#3 0.2.1.1.1.1 (2+3) C+ ConditionsFromPly = 4 Anti-Andernach Disparate

#### 19 - Scheme



#2 (7+3) C+ ConditionsUntilPly = 1 Anti-Andernach

n°20. Here after the A fairy part (two half-moves) follows the B orthodox part (only one half-move which is a mating move).

Because every #2 has only three half-moves (plies), it means that the 1<sup>st</sup> ply (white key-move) and 2<sup>nd</sup> ply (black defence) are under the condition Anti-Andernach, followed by the last 3<sup>rd</sup> orthodox ply (white mating move).

It is obvious that here we have to start (as the  $1^{st}$  ply) with the white Pawn b7, which can immediately promote into a black figure! The black defence ( $2^{nd}$  ply) is a move with this black piece that instantly turns into a white figure. So ends the first fairy phase of the problem, as the condition Anti-Andernach works only to this point (until ply = 2). Next we must give mate with the new white figure on the board.

Of course : 1.b8 = #(#)? blocus is not a good idea, for example : one refutation  $1... \textcircled{#} \times 65 + !$ 

But another attempt:  $1.b8 = \mathbb{Z}(\mathbb{Z})$ ? blocus deserves attention! Now Black is in zugzwang and can only play with his Rook. But it's interesting that this figure has a lot of opportunities for defence:

1...  $\square$  b8( $\square$ )~ (along the 8<sup>th</sup> rank) ends the Anti-Andernach phase and is followed by the orthodox move : 2.  $\square$  d8#

But there are here many additional black defences that show «Black correction»:

- 1...**ਡ**b7(**ਡ**)! 2.**ਡ**d7#
- 1...**፮**b6(፮)! 2.፮d6#
- 1... **≝**b5( **≌**)! 2.ç6#
- 1... **■** b4( **■**)! 2. **■** d4#
- 1... **■**b2(**□**)! 2. **□**d2#
- 1...**ਡ**b1(**ਡ**)! 2.**ਡ**d1#

But  $1... \blacksquare b3(\blacksquare)!!$  and no mate.

Correct is  $1.b8 = \emptyset$  (4)! blocus and now follows:

- 1...**△**d7(⑤) 2.⑤b6#
- 1... **2 c**6(**2**) **2**.**2 e**7#
- 1...**♠** a6(♠) 2.♠ç7#

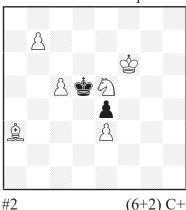
Of course, using this matrix you can compose a lot of direct problems with similar contents with multiple promotions of white Pawns(s)!

 $n^{\circ}21$ : As we have already said, each problem with stipulation sh# (=, ==, etc...), ss# (=, ==, etc...) etc... in n moves, has n + 1 half-moves (plies), since all the moves (except the last, combined move of both sides) are single moves.

In other words, here we have in total seven half-moves, five of them under the condition of Sentinels. After this follow two orthodox half-moves. Here's how it works:

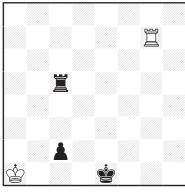
1.  $\square g2(+ \triangle g7)$  2. g8= 2 3 3. a2 4 4.  $b2(+ \triangle a2)$  5.  $x \approx c2 (+ \triangle b2)$  This is the end of the fairy-part with Sentinels condition. Now follows the orthodox finale: 6.  $x \approx c1 + 2 \times c1 = c1$ 

**20 - P. Petkov** Education example



ConditionsUntilPly = 2
Anti-Andernach

#### 21 - Scheme



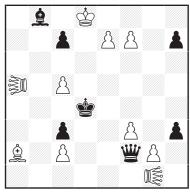
ss#6 (2+3) C+ ConditionsUntilPly = 5 Sentinels The 22 is the best problem of all those I created using WinChloe's options! Here ConditionsUntilPly=4, means that condition Anti-Andernach is valid only for the first four half-moves (plies) i.e. until the second move of White and Black. Then the remaining two halfmoves (the 3<sup>rd</sup> move of White and Black) are performed as orthodox half-moves!

Some explanations about the solutions: here alphabetic symbols (X, Y, Z for the key-moves and A, B, C for the second white moves) are needed to make easier to understand the cyclical motives in anti-dual attempts at the formation of black batteries.

#### Solutions:

- a) 1.é8= ②(全)! X 豐×f3! 2.f8= 豐(豐)! A 豐c6(豐) 3.豐a4+ 全×a4#  $[2.f8 = \mathbb{Z}(\mathbb{Z})? \mathbf{B}, 2.f8 = \mathbb{Q}(\mathbb{Q})? \mathbf{C}]$
- b) 1.é8=公(**1**)! Y 豐×ç2! 2.f8=罩(罩)! B 豐f5(豐) 3.豐d7+ **1**d6# [2.f8=隱(**低**)? C, 2.f8=豐(**豐**)? A]
- c) 1.é8=劉(劉)! Z 豐×g2! 2.f8=奚(溪) C 豐g8(豐) 3.豐h8+ 劉×h8#

22 - P. Petkov KoBulChess 2017 Dedicated to Christian Poisson



hs#3 (10+7) C+Conditions Until Ply = 4b) \(\displaystyle{d} d8 \rightarrow \cong 8 c) **3**a5→d7

Anti-Andernach

**I**€=Locuste **≥**Lion

#### Thematic complex:

- 1 Super AUW, realised in three phases after cyclic six-fold promotions of the white Pawns é7 and f7, that form black batteries! This idea is combined with a double cyclic anti-dual after promotions of the rear battery pieces! This idea and mechanism are shown for the first time! Here's the cyclic formula:
  - a) 1.X! 2.A! B? C?
  - b) **1.Y!** 2.B! C? A?
  - c) 1.Z! 2.C! A? B?
- 2 The black moves: a) 1... \subseteq \times f3!; b) 1... \subseteq \times c2!; c) 1... \subseteq \times g2! aim to keep the black colour for the black Queen; to open a line for the white Locust.
- 3 During the play the black Queen makes a record number of moves a total of 9 in the three phases! I think that this form of Super AUW - with six different promotions to black figures from only two white Pawns (in addition with thematic cyclical anti-dual motives connected with the promotions of the rear figures of the black batteries!), is shown here for the first time in a fairy chess composition!

#### 3. Combined play with both options

This is perhaps the most complex and interesting combination of options. Unfortunately, probably there are no published examples of this type so far. So we will briefly explain the definition with a light scheme.

For example, if the total number of half-moves is X, we can note:

X = A orthodox half-moves + B fairy half-moves + C orthodox half-moves (combining ConditionsFromPly and ConditionsUntilPly).

If we have a problem with total X = 10 half-moves, a possible variant can be the following:

X = 3 orthodox half-moves + 4 fairy half-moves + 3 orthodox half-moves.

If we use the condition «Circe» below the diagram in the twin field of WinChloe we must write:

ConditionsFromply=4

ConditionsUntilPly=7

And in the conditions field we must of course mark «Circe».

In other words here our condition Circe «works» only between the 4th ply and 7th ply. All other halfmoves (from 1st to 3rd and from 8th to 10th) are orthodox plies!

Of course, it is not easy to compose sufficiently interesting problems with such a unique «formula» because it is rather difficult to motivate the necessity (expediency) from such a difficult structure of half-moves (plies). I think that drawing up such tasks is to a large extent a test for the composer's fantasy, or more precisely to say, a test of his ability for a non-standard mission!

Now follows a light example. Here we have : two orthodox half-moves (plies) + three fairy (Sentinels) half-moves (plies) + two last orthodox half-moves (plies).

#### $n^{\circ}23 : 1. \Xi h2 2. \Xi a2$

End of first orthodox part which was intended to prepare for the next «blockade cannonade» of the condition Sentinels:

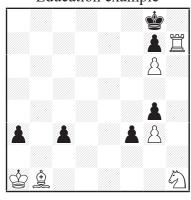
3.  $\square$   $\varsigma$ 2(+  $\triangle$  a2) 4.  $\square$  f2(+  $\triangle$   $\varsigma$ 2) 5.  $\square$  h2(+  $\triangle$  f2)

End of fairy part. Now follows the second orthodox part:

#### 6. □ h8+ •×h8=

Pay attention to the very important feature of the final part : after 6. \( \mathbb{\subset} \) h8+ the Sentinels condition no longer applies and a white Pawn is not put on the square h2!

**23 - P. Petkov** Education example



ss=6

(6+6) C+

ConditionsFromPly = 3 ConditionsUntilPly = 5

Sentinels

#### III. CONCLUSION

In my opinion, almost certainly problems of WinChloe's options have a great future. Of course, in this respect very thorough and comprehensive work is needed. Many tests are needed for the various features offered by the options. I am convinced that play with change of conditions within a single solution or single variant is a very avant-garde approach in Fairy Chess, although almost similar ideas have been defined and slightly experimented with in the past.

It would be very interesting to know if Christian is going to work in the future in this arena, since his options are certainly a very rich opportunity for a new discovery on the same basis !? I have a positive expectation in this regard !

Speaking of these non-standard forms, I would like to briefly mention that in the year 1999, in the German magazine, *feenschach*, (Heft 134, Band XXVI, November-December 1999) was published my article: «Problems and endgames with change of conditions (BW)». BW in German means BedingungsWechsel (change of conditions). In this long article (nine pages) dedicated to the 50<sup>th</sup> anniversary of the magazine, I offered different versions of BW, but my approach was quite different - a change of the fairy condition on every move by both sides (change after every two half-moves). Since I'm not a programmer, I cannot say at all whether it was possible to program my «BW» ideas. On this question, I cannot give an answer to this day...

Thus, the idea of BW is still not very popular and I have composed few originals of this type (some of them can be found in the databases of WinChloe).

But I must repeat: Christian Poisson's approach to solving the problem about change of fairy conditions is fundamentally new and already works well! So the name of such problems WinChloe's options seems to me quite natural!

(■ Petko Petkov, September 2018, with the help of Geoff Foster for translation in english)